INFORMATION THEORY & CODING Week 10 : Differential Entropy

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- Definitions
- AEP for Continuous Random Variables
- Relation of differential entropy to discrete entropy



From discrete to continuous variables



Probability Distribution of X





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Definition

Let X be a random variable with cumulative distribution function (CDF) $F(x) = \Pr(X \le x)$. If F(x) is continuous, the random variable is continuous. Let f(x) = F'(X) when the derivative is defined. If $\int_{-\infty}^{+\infty} f(x) = 1$, f(x) is called the probability density function (pdf) for X. The set of x where f(x) > 0 is called the support set of the X.

Definition

The differential entropy $h(\boldsymbol{X})$ of a continuous random variable \boldsymbol{X} with density $f(\boldsymbol{x})$ is defined as

$$h(X) = -\int_{\mathcal{S}} f(x) \log f(x) dx = h(f),$$

where \mathcal{S} is the support set of the random variable.

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•
$$f(x) = \frac{1}{a}, x \in [0, a]$$

• The differential entropy is:

$$h(X) = -\int_0^a \frac{1}{a} \log \frac{1}{a} \mathrm{d}x = \log a$$
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• for a < 1, $h(X) = \log a < 0$, differential entropy can be negative! (unlike discrete entropy)



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Example: Normal distribution

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$$X \sim \phi(x) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp(\frac{-x^2}{2\sigma^2}), x \in \mathbb{R}$$

• Differential entropy:

$$h(\phi) = rac{1}{2} \log 2\pi e \sigma^2$$
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Calculation:

$$h(\phi) = -\int \phi \log \phi dx = -\int \phi(x) \left[-\frac{x^2}{2\sigma^2} \log e - \log \sqrt{2\pi\sigma^2} \right] dx$$
$$= \frac{\mathbb{E}(X^2)}{2\sigma^2} \log e + \frac{1}{2} \log 2\pi\sigma^2 = \frac{1}{2} \log e + \frac{1}{2} \log 2\pi\sigma^2$$
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• Discrete world: for a sequence of i.i.d. random variables

$$\frac{1}{n}\log p(X_1, X_2, \dots, X_n) \to H(X).$$

• Continuous world: for a sequence of i.i.d. random variables

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Proof follows from the weak law of large numbers.



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Typical set

• Discrete case: number of typical sequences

 $\left|A_{\epsilon}^{(n)}\right| \approx 2^{nH(X)}$

• Continuous case: The volume of the typical set

$$\operatorname{Vol}(A) = \int_A \mathrm{d}x_1 \mathrm{d}x_2 \dots \mathrm{d}x_n, \ A \subset \mathbb{R}^n.$$

Definition

For $\epsilon > 0$ and any n, we define the typical set $A_{\epsilon}^{(n)}$ with respect to f(x) as follows:

$$A_{\epsilon}^{(n)} = \left\{ (x_1, x_2, \dots, x_n) \in \mathcal{S}^n : \left| -\frac{1}{n} \log f(x_1, x_2, \dots, x_n) - h(X) \right| \le \epsilon \right\},$$

where $f(x_1, x_2, ..., x_n) = \prod_{i=1}^n f(x_i)$.

The typical set $A_{\epsilon}^{(n)}$ has the following properties: 1. $\Pr(A_{\epsilon}^{(n)}) > 1 - \epsilon$ for n sufficiently large. 2. $\operatorname{Vol}(A_{\epsilon}^{(n)}) \leq 2^{n(h(X)+\epsilon)}$ for all n. 3. $\operatorname{Vol}(A_{\epsilon}^{(n)}) \geq (1-\epsilon)2^{n(h(X)-\epsilon)}$ for n sufficiently large.

Proof. 1.

Similar to the discrete case. By definition, $-\frac{1}{n}\log f(X^n) = -\frac{1}{n}\sum \log f(X_i) \rightarrow h(X)$ ir probability.



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The typical set $A_{\epsilon}^{(n)}$ has the following properties: 2. $\operatorname{Vol}(A_{\epsilon}^{(n)}) \leq 2^{n(h(X)+\epsilon)}$ for all n.

Poof. 2.

$$1 = \int_{\mathcal{S}^n} f(x_1, x_2, \dots, x_n) dx_1 dx_2 \dots dx_n$$

$$\geq \int_{A_{\epsilon}^{(n)}} f(x_1, x_2, \dots, x_n) dx_1 dx_2 \dots dx_n$$

$$\geq \int_{A_{\epsilon}^{(n)}} 2^{-n(h(X)+\epsilon)} dx_1 dx_2 \dots dx_n = 2^{-n(h(X)+\epsilon)} \int_{A_{\epsilon}^{(n)}} dx_1 dx_2 \dots dx_n$$

$$= 2^{-n(h(X)+\epsilon)} \mathsf{Vol}(A_{\epsilon}^{(n)}).$$

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The typical set $A_{\epsilon}^{(n)}$ has the following properties: 3. $\operatorname{Vol}(A_{\epsilon}^{(n)}) \geq (1-\epsilon)2^{n(h(X)-\epsilon)}$ for n sufficiently large.

Proof. 3.

$$1 - \epsilon \leq \int_{A_{\epsilon}^{(n)}} f(x_1, x_2, \dots, x_n) dx_1 dx_2 \dots dx_n$$
$$\leq \int_{A_{\epsilon}^{(n)}} 2^{-n(h(X)-\epsilon)} dx_1 dx_2 \dots dx_n$$
$$= 2^{-n(h(X)-\epsilon)} \int_{A_{\epsilon}^{(n)}} dx_1 dx_2 \dots dx_n$$
$$= 2^{-n(h(X)-\epsilon)} \operatorname{Vol}(A_{\epsilon}^{(n)}).$$

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An interpretation

- The volume of the smallest set that contains most of the probability is approximately $2^{nh(X)}$.
- For an *n*-dim volume, this means that each dim has length $(2^{nh(X)})^{\frac{1}{n}} = 2^{h(X)}$.



Mean value theorem (MVT)

If a function f is continuous on the closed interval [a, b], and differentiable on (a, b), then there exists a point $c \in (a, b)$ such that



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- Consider a random variable X with pdf f(x). We divide the range of X into bins of length Δ.
- MVT: there exists a value $x_i \in (i\Delta, (i+1)\Delta)$ within each bin such that



• Define the quantized random variable as $X^{\Delta} = x_i$ if $i\Delta \leq X \leq (i+1)\Delta$ with pmf

$$p_i = \Pr[X^{\Delta} = x_i] = \int_{i\Delta}^{(i+1)\Delta} f(x) \mathrm{d}x = f(x_i)\Delta.$$

• The entropy of
$$X^{\Delta}$$
 is

$$H(X^{\Delta}) = -\sum_{-\infty}^{+\infty} p_i \log p_i = -\sum_{-\infty} \Delta f(x_i) \log f(x_i) - \log \Delta.$$

• If f(x) is is Riemann integrable, as $\Delta \rightarrow 0$,

 $H(X^{\Delta}) + \log \Delta \to h(f) = h(X)$



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$$H(X^{\Delta}) + \log \Delta \to h(f) = h(X)$$

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- Reading: Chapter 8: 8.1 8.3
- Homework: Problems 8.1, 8.5, 8.7

